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# Cryptography I

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General concepts and some classical  
ciphers

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# Security Goals

- Confidentiality (secrecy, privacy)
    - Assure that data is accessible to only one who are authorized to know
  - Integrity
    - Assure that data is only modified by authorized parties and in authorized ways
  - Availability
    - Assure that resource is available for authorized users
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# General tools

- Cryptography
  - Software controls
  - Hardware controls
  - Policies and procedures
  - Physical controls
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# What is Crypto?

- Constructing and analyzing **cryptographic protocols** which enable parties to achieve security objectives
    - Under the presence of adversaries.
  - A protocol (or a scheme) is a suite of procedures that tell each party what to do
    - usually, computer algorithms
  - Cryptographers devise and analyze protocols under **Attack model**
    - assumptions about the resources and actions available to the adversary
      - So, you need to think as an adversary
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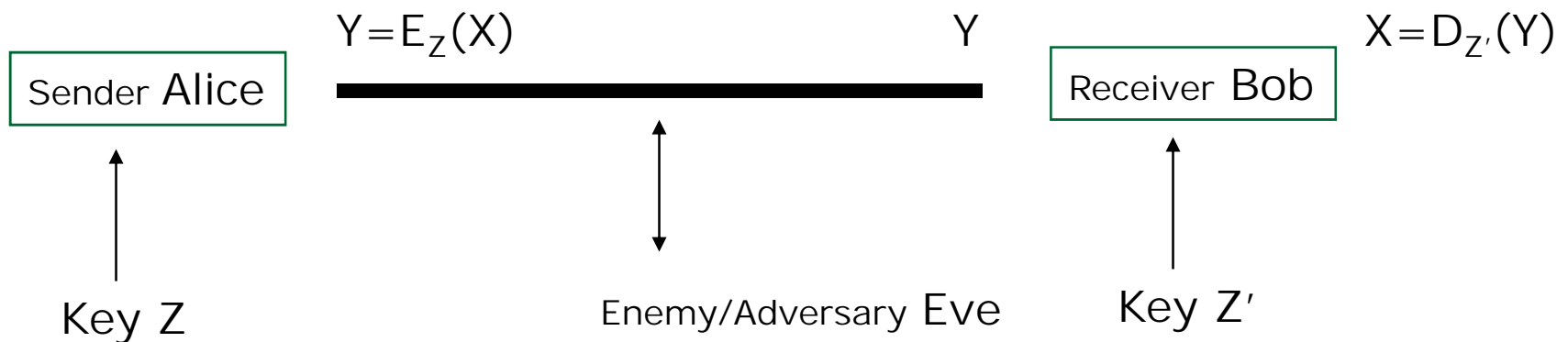
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# Terms

- **Cryptography:** the study of mathematical techniques for providing information security services.
  - **Cryptanalysis:** the study of mathematical techniques for attempting to get security services breakdown.
  - **Cryptology:** the study of cryptography and cryptanalysis.
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# Terms ...

- plaintexts
- ciphertexts
- keys
- encryption
- decryption



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# Secret-key cryptography

- Also called: symmetric cryptography
  - Use the same key for both encryption & decryption ( $Z=Z'$ )
  - Key must be kept secret
  - Key distribution – how to share a secret between A and B very difficult
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# Public-key cryptography

- Also called: asymmetric cryptography
  - Encryption key different from decryption key and
    - It is not possible to derive decryption key from encryption key
  - Higher cost than symmetric cryptography
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# Is it a secure cipher system?

- **Why insecure**

- **just break it under a certain reasonable attack model (show failures to assure security goals)**

- **Why secure:**

- Evaluate/prove that under the considered attack model, security goals are assured
  - Provable security: Formally show that (with mathematical techniques) the system is as secure as a well-known secure one (usually simpler).
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# Breaking ciphers ...

- There are different methods of breaking a cipher, depending on:
    - the type of information available to the attacker
    - the interaction with the cipher machine
    - the computational power available to the attacker
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# Breaking ciphers ...

## ■ **Ciphertext-only attack:**

- The cryptanalyst knows **only the ciphertext**.
- Goal: to find the plaintext and the key.
- NOTE: such vulnerable is seen completely insecure

## ■ **Known-plaintext attack:**

- The cryptanalyst knows **one or several pairs of ciphertext and the corresponding plaintext**.
  - Goal: to find the key used to encrypt these messages
    - or a way to decrypt any new messages that use the same key (although may not know the key).
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# Breaking ciphers ...

- **Chosen-plaintext attack**

- The cryptanalyst **can choose a number of messages and obtain the ciphertexts for them**
- Goal: deduce the key used in the other encrypted messages or decrypt any new messages (using that key).

- **Chosen-ciphertext attack**

- Similar to above, but the cryptanalyst **can choose a number of ciphertexts and obtain the plaintexts.**

- Both can be **adaptive**

- The choice of ciphertext may depend on the plaintext received from previous requests.
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# Models for Evaluating Security

- **Unconditional (information-theoretic) security**
    - **Assumes that the adversary has unlimited computational resources.**
    - Plaintext and ciphertext modeled by their distribution
    - Analysis is made by using probability theory.
    - For encryption systems: **perfect secrecy**, observation of the ciphertext provides no information to an adversary.
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# Models for Evaluating Security

- **Provable security:**

- Prove security properties based on assumptions that it is difficult to solve a well-known and supposedly difficult problem (NP-hard ...)
  - E.g.: computation of discrete logarithms, factoring

- **Computational security (practical security)**

- Measures the amount of computational effort required to defeat a system using the best-known attacks.
  - Sometimes related to the hard problems, but no proof of equivalence is known.
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# Models for Evaluating Security

- **Ad hoc security (heuristic security):**
    - Variety of convincing arguments that every successful attack requires more resources than the ones available to an attacker.
    - Unforeseen attacks remain a threat.
    - **THIS IS NOT A PROOF**
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# **Classic ciphers**

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# Shift cipher (additive cipher)

- Key Space: [1 .. 25]
- Encryption given a key K:
  - each letter in the plaintext P is replaced with the K'th letter following corresponding number (shift right):
  - Another way:  $Y = X \oplus K \rightarrow$  additive cipher
- Decryption given K:
  - shift left

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

P = CRYPTOGRAPHYISFUN

K = 11

C = NCJAVZRCLASJTDQFY

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# Shift Cipher: Cryptanalysis

- Easy, just do exhaustive search
  - key space is small ( $\leq 26$  possible keys).
  - once  $K$  is found, very easy to decrypt



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# General Mono-alphabetical Substitution Cipher

- The key space: all permutations of  $\Sigma = \{A, B, C, \dots, Z\}$
- Encryption given a key  $\pi$ :
  - each letter  $X$  in the plaintext  $P$  is replaced with  $\pi(X)$
- Decryption given a key  $\pi$  :
  - each letter  $Y$  in the ciphertext  $P$  is replaced with  $\pi^{-1}(Y)$

- **Example:**

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z  
 $\pi =$  B A D C Z H W Y G O Q X S V T R N M S K J I P F E U

BECAUSE  $\rightarrow$  AZDBJSZ

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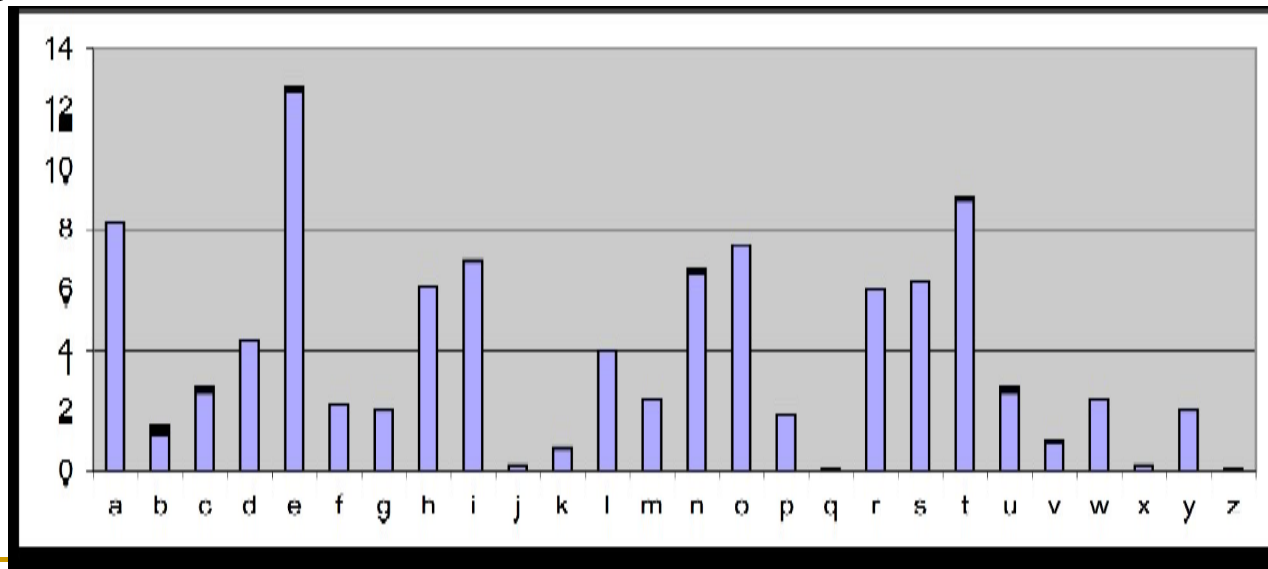
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# Looks secure, early days

- Exhaustive search is infeasible
    - key space size is  $26! \approx 4 \cdot 10^{26}$
  - Dominates the art of secret writing throughout the first millennium A.D.
  - Thought to be unbreakable by many back then
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# Cryptanalysis of Substitution Ciphers: Frequency Analysis

- Each language has certain features:
  - frequency of letters, or of groups of two or more letters.
- Substitution ciphers preserve the mentioned language features → vulnerable to frequency analysis attacks



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# Substitution Ciphers: Cryptanalysis

- The number of different ciphertext characters or combinations are counted to determine the frequency of usage.
- The cipher text is examined for patterns, repeated series, and common combinations.
- Replace ciphertext characters with possible plaintext equivalents using known language characteristics.

- Example:

THIS IS A PROPER SAMPLE FOR ENGLISH TEXT. THE FREQUENCIES OF LETTERS IN THIS SAMPLE IS NOT UNIFORM AND VARY FOR DIFFERENT CHARACTERS. IN GENERAL THE MOST FREQUENT LETTER IS FOLLOWED BY A SECOND GROUP. IF WE TAKE A CLOSER LOOK WE WILL NOTICE THAT FOR BIGRAMS AND TRIGRAMS THE NONUNIFORM IS EVEN MORE.

- Observations:  $f_x=1$  và  $f_A=15$ .
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- The letters in the English alphabet can be divided into 5 groups of similar frequencies

I e

II t,a,o,i,n,s,h,r

III d,l

VI c,u,m,w,f,g,y,p,b

V v,k,j,x,q,z

- Some frequently appearing bigrams or trigrams

Th, he, in, an, re, ed, on, es, st, en at, to

The, ing, and, hex, ent, tha, nth, was eth, for, dth.

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# Example

Letter:	A	B	C	D	E	F	G
Frequency:	5	24	19	23	12	7	0
Letter:	H	I	J	K	L	M	N
Frequency:	24	21	29	6	21	1	3
Letter:	O	P	Q	R	S	T	U
Frequency:	0	3	1	11	14	8	0
Letter:	V	W	X	Y	Z		
Frequency:	27	5	17	12	45		

■  $e \Rightarrow Z$

$f_j = 29, f_v = 27$

$f_{jcz} = 8 \rightarrow t \Rightarrow J$

$h \Rightarrow C$

■  $a \Rightarrow V$

(article a)

$J, V, B, H, D, I, L, C \{t, a, o, i, n, s, h, r\}$

$t, a$                        $h$

$JZB = te ? \{teo, tei, ten, ter, tes\} \rightarrow n \Rightarrow B$



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- Observations:

- A cipher system should not allow statistical properties of plaintext to pass to the ciphertext.
- The ciphertext generated by a "good" cipher system should be statistically indistinguishable from random text.

- Idea for a stronger cipher (1460's by Alberti)

- use more than one cipher alphabet, and switch between them when encrypting different letters → Polyalphabetic Substitution Ciphers
  - Developed into a practical cipher by Vigenère (published in 1586)
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- **Definition:**

- Given  $m$ , a positive integer,  $P = C = (\mathbb{Z}_{26})^n$ , and  $K = (k_1, k_2, \dots, k_m)$  a key, we define:

- **Encryption:**

$$e_k(p_1, p_2 \dots p_m) = (p_1+k_1, p_2+k_2 \dots p_m+k_m) \pmod{26}$$

- **Decryption:**

$$d_k(c_1, c_2 \dots c_m) = (c_1-k_1, c_2-k_2 \dots c_m-k_m) \pmod{26}$$

- **Example:**

Plaintext: C R Y P T O G R A P H Y

Key: L U C K L U C K L U C K

Ciphertext: N L A Z E I I B L J J I

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# Vigenere Cipher: Cryptanalysis

- Find the length of the key.
  - Divide the message into that many shift cipher encryptions.
  - Use frequency analysis to solve the resulting shift ciphers.
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# One-Time Pad

Key is chosen randomly

Plaintext  $X = (x_1 \ x_2 \ \dots \ x_n)$

Key  $K = (k_1 \ k_2 \ \dots \ k_n)$

Ciphertext  $Y = (y_1 \ y_2 \ \dots \ y_n)$

$$e_k(X) = (x_1+k_1 \ x_2+k_2 \ \dots \ x_n+k_n) \bmod m$$

$$d_k(Y) = (x_1-k_1 \ x_2-k_2 \ \dots \ x_n-k_n) \bmod m$$

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# Example

Plaintext space = Ciphertext space =

Keyspace =  $\{0,1\}^n$

Key is chosen randomly

For example:

Plaintext is                   10001011

Key is                           00111001

Then ciphertext is           10110010

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# Main points in One-Time Pad

- The key is never to be reused
    - Thrown away after first and only use
    - If reused → insecure!
  - One-Time Pad uses a very long key, exactly the same length as of the plaintext
    - In old days, some suggest choose the key as texts from, e.g., a book → i.e. not **randomly chosen**
      - Not One-Time Pad anymore → this does not have perfect secrecy as in true One-Time-Pad and can be broken
    - Perfect secrecy means key length be at least message length
      - **Difficult in practice!**
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- Shift ciphers are easy to break using brute force attacks (exhaustive key search)
  - Substitution ciphers preserve language features (in N-gram frequency) and are vulnerable to frequency analysis attacks.
  - Vigenère cipher are also vulnerable to frequency analysis once the key length is found.
    - In general poly-alphabetical substitution ciphers are not that secure
  - OTP has perfect secrecy if the key is chosen randomly in the message length and is used only once.
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