

Data compression

anhtt-fit@mail.hut.edu.vn

dungct@it-hut.edu.vn

Data Compression

- Data in memory have used fixed length for representation
- For data transfer (in particular), this method is inefficient.
- For speed and storage efficiencies, data symbols should use the minimum number of bits possible for representation.
- Methods Used For Compression:
 - Encode high probability symbols with fewer bits
 - Shannon-Fano, Huffman, UNIX compact
 - Encode sequences of symbols with location of sequence in a dictionary
 - PKZIP, ARC, GIF, UNIX compress, V.42bis
 - Lossy compression
 - JPEG and MPEG

Variable Length Bit Codings

- Suppose 'A' appears 50 times in text, but 'B' appears only 10 times
- ASCII coding assigns 8 bits per character, so total bits for 'A' and 'B' is $60 * 8 = 480$
- If 'A' gets a 4-bit code and 'B' gets a 12-bit code, total is $50 * 4 + 10 * 12 = 320$

Compression rules:

- Use minimum number of bits
- No code is the prefix of another code
- Enables left-to-right, unambiguous decoding

Variable Length Bit Codings

- No code is a prefix of another
 - For example, can't have 'A' map to 10 and 'B' map to 100, because 10 is a prefix (the start of) 100.
- Enables left-to-right, unambiguous decoding
 - That is, if you see 10, you know it's 'A', not the start of another character.

Variable-length encoding

- Use different number of bits to encode different characters.
- Ex. Morse code.
- Issue: ambiguity.

- • • • - - - • • •
- SOS ?
- IAMIE ?
- EEWNI ?
- V7O ?

S	•••	← prefix of V
E	•	← prefix of I, S
I	••	← prefix of S
V	•••-	

Letters	Numbers
A	1
B	2
C	3
D	4
E	5
F	6
G	7
H	8
I	9
J	0
K	
L	
M	
N	
O	
P	
Q	
R	
S	
T	
U	
V	
W	
X	
Y	
Z	

Huffman code

- Constructed by using a code tree, but starting at the leaves
- A compact code constructed using the binary Huffman code construction method

Huffman code Algorithm

- ① **Make a leaf node for each code symbol**
 - Add the generation probability of each symbol to the leaf node
- ② **Take the two leaf nodes with the smallest probability and connect them into a new node**
 - Add 1 or 0 to each of the two branches
 - The probability of the new node is the sum of the probabilities of the two connecting nodes
- ③ **If there is only one node left, the code construction is completed. If not, go back to (2)**

Demo

- [65demo-huffman.ppt](#)

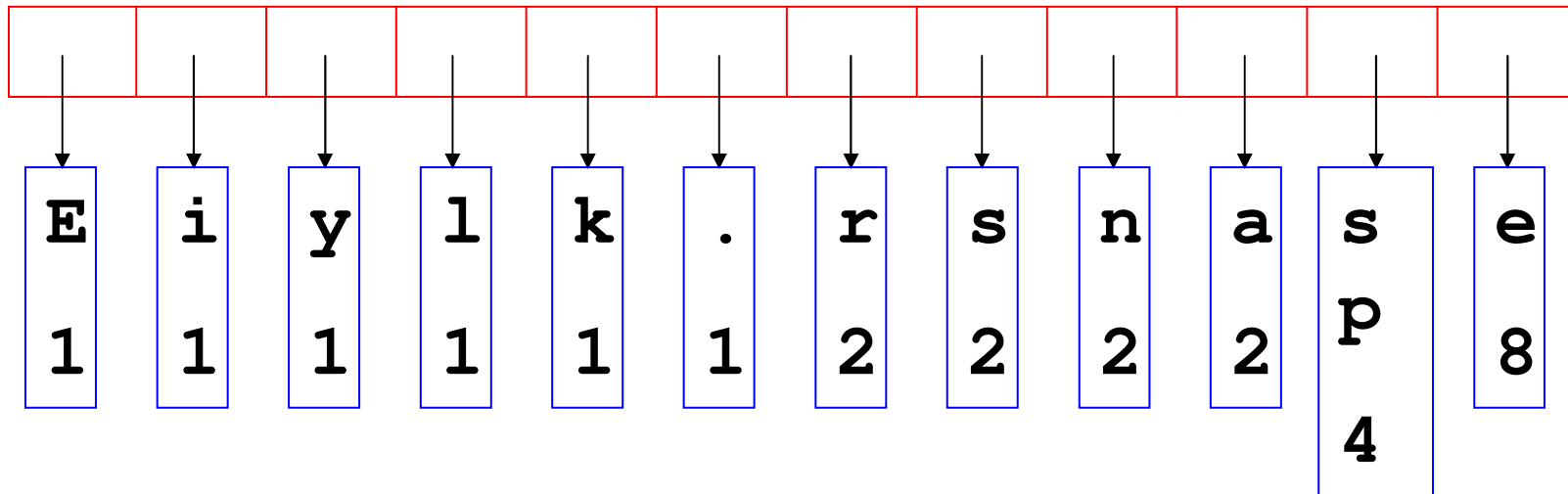
Compress a text

- Consider the following short text:
Eerie eyes seen near lake.
- Count up the occurrences of all characters in the text

Char	Freq.	Char	Freq.	Char	Freq.
E	1	y	1	k	1
e	8	s	2	.	1
r	2	n	2		
i	1	a	2		
space	4	l	1		

Building a Tree

- The queue after inserting all nodes

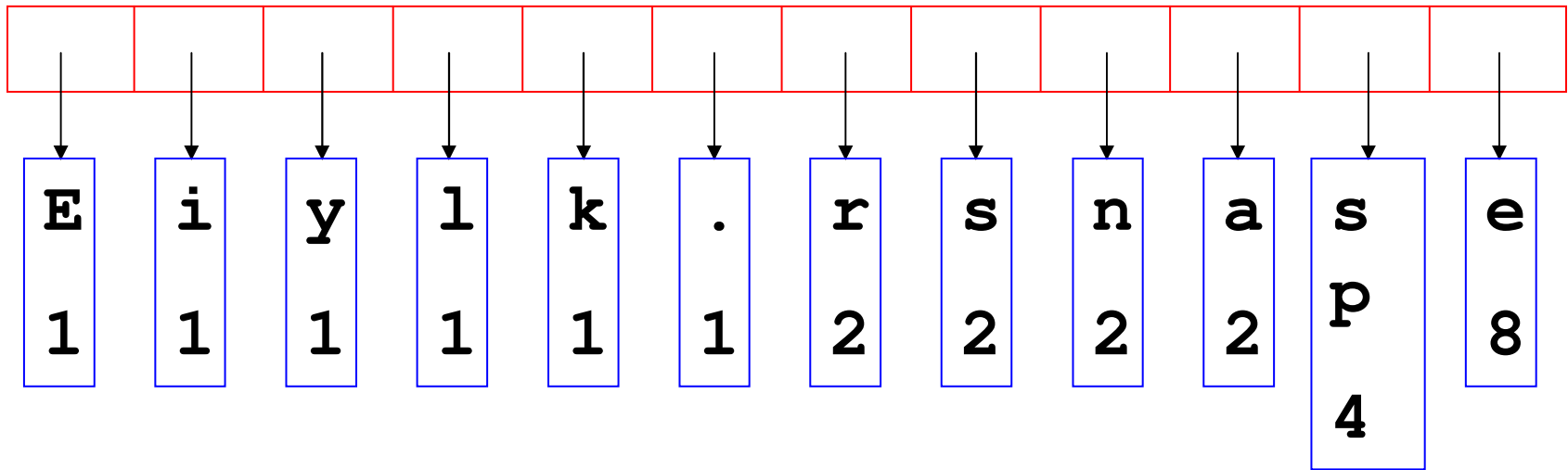


- Null Pointers are not shown

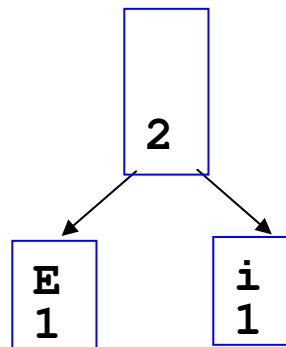
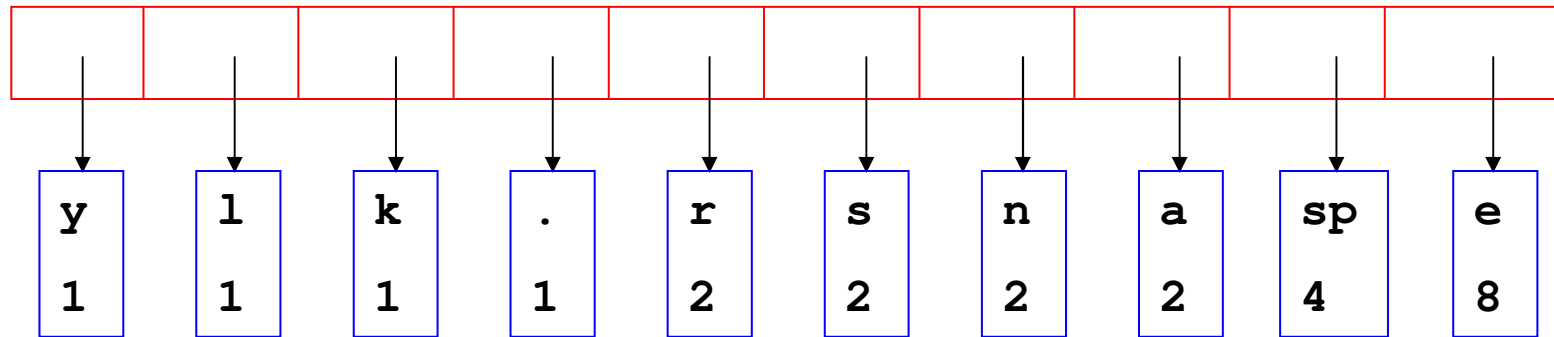
Building a Tree

- While priority queue contains two or more nodes
 - Create new node
 - Dequeue node and make it left subtree
 - Dequeue next node and make it right subtree
 - Frequency of new node equals sum of frequency of left and right children
 - Enqueue new node back into queue

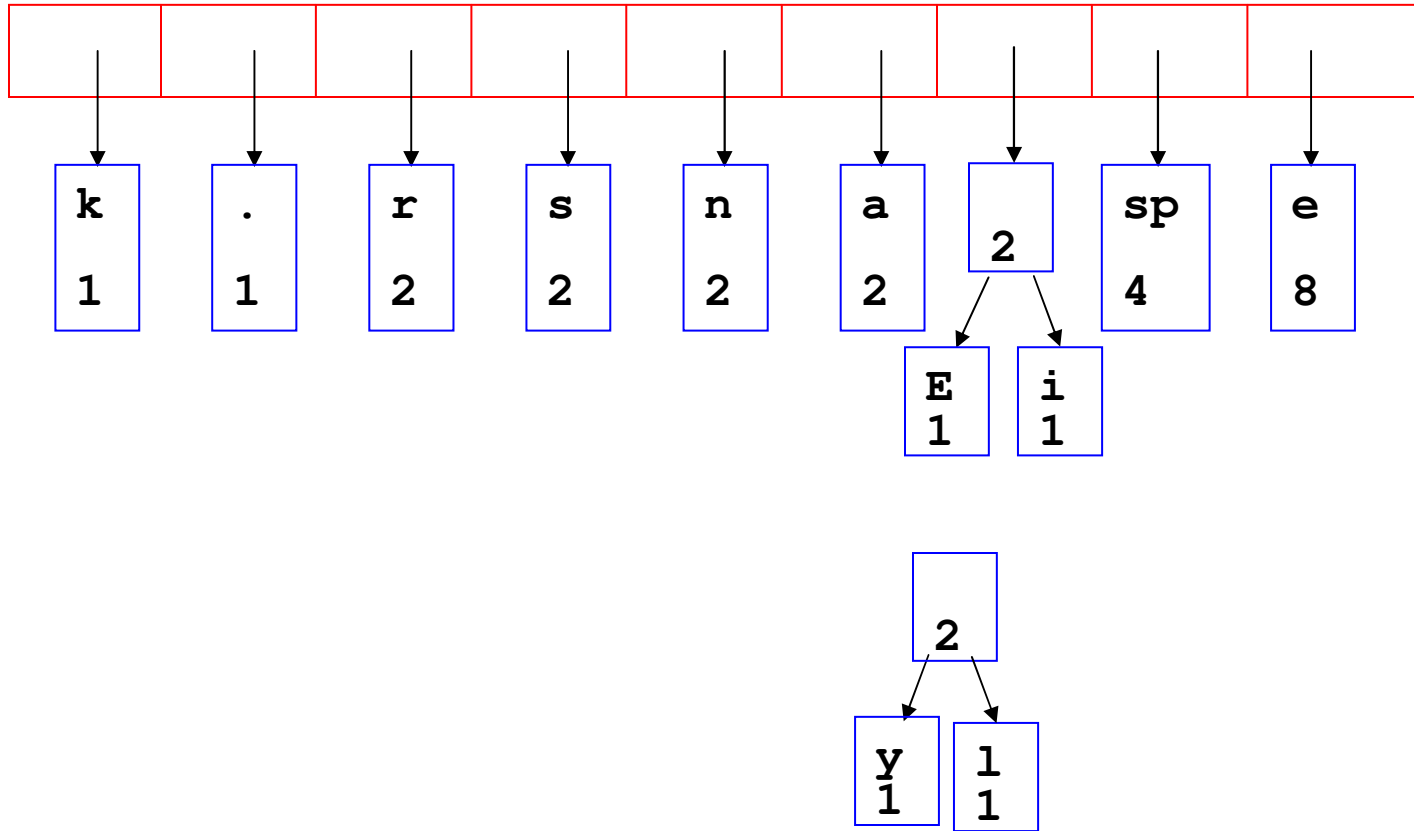
Building a Tree



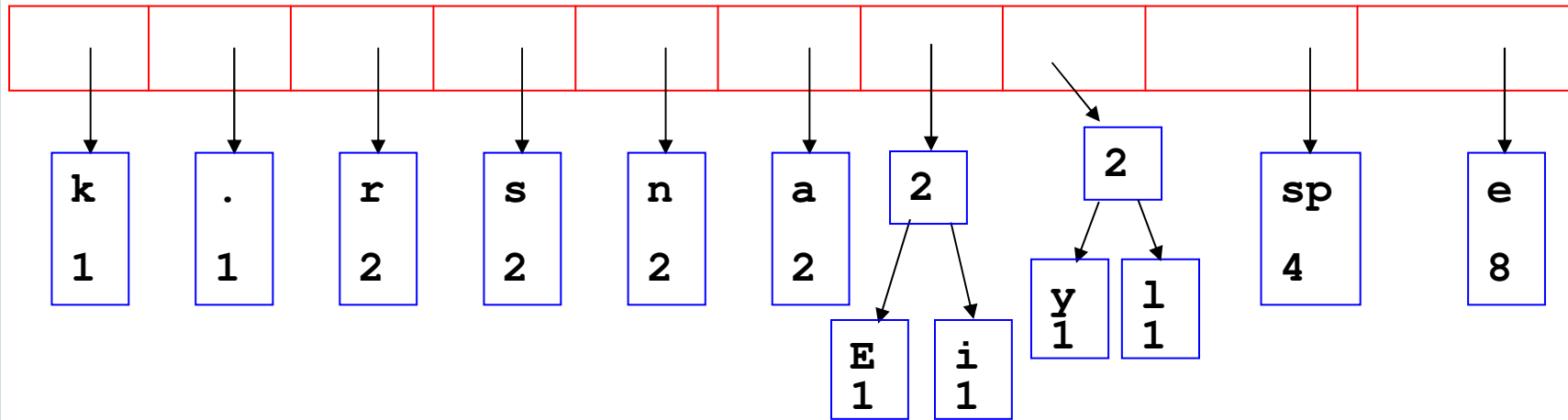
Building a Tree



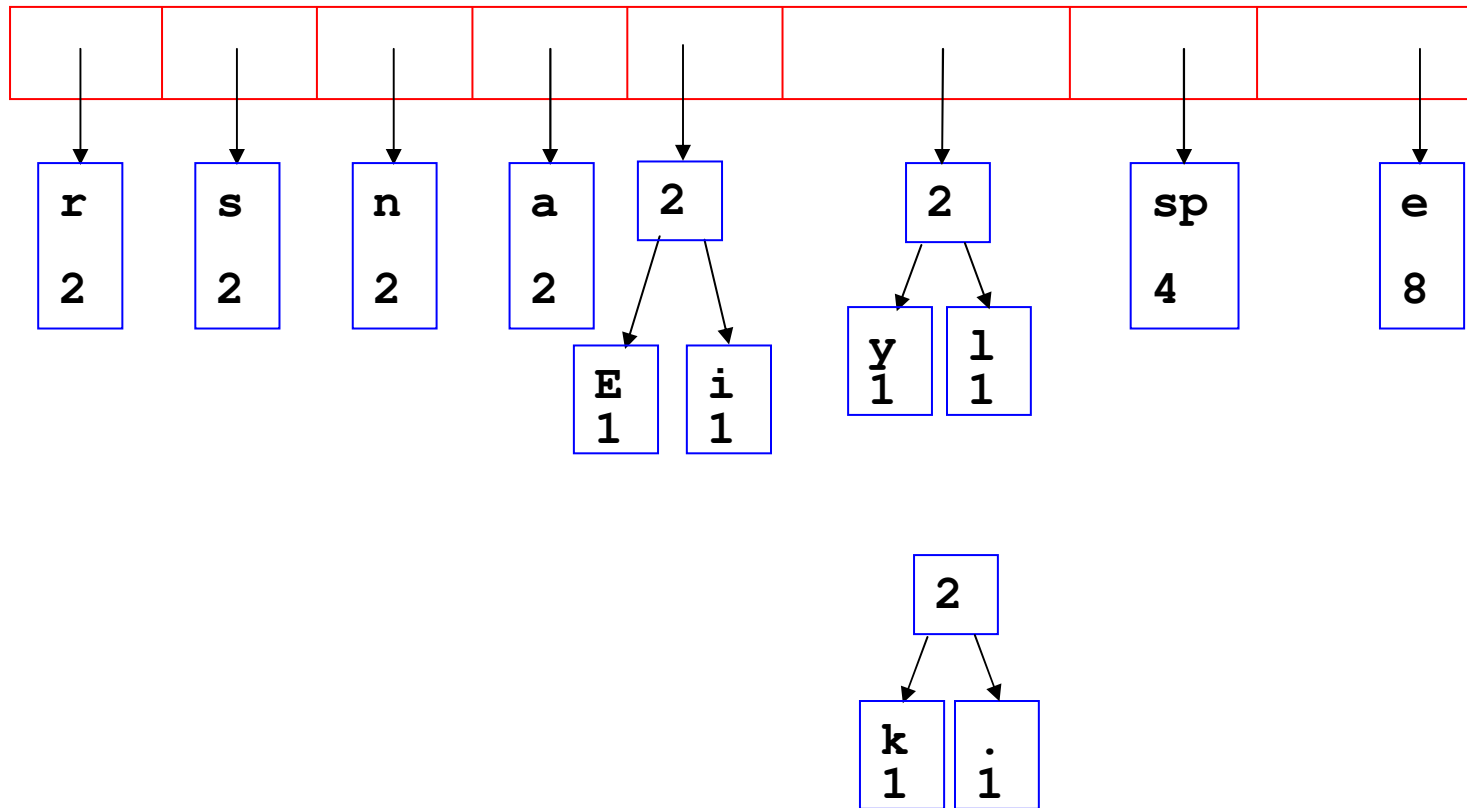
Building a Tree



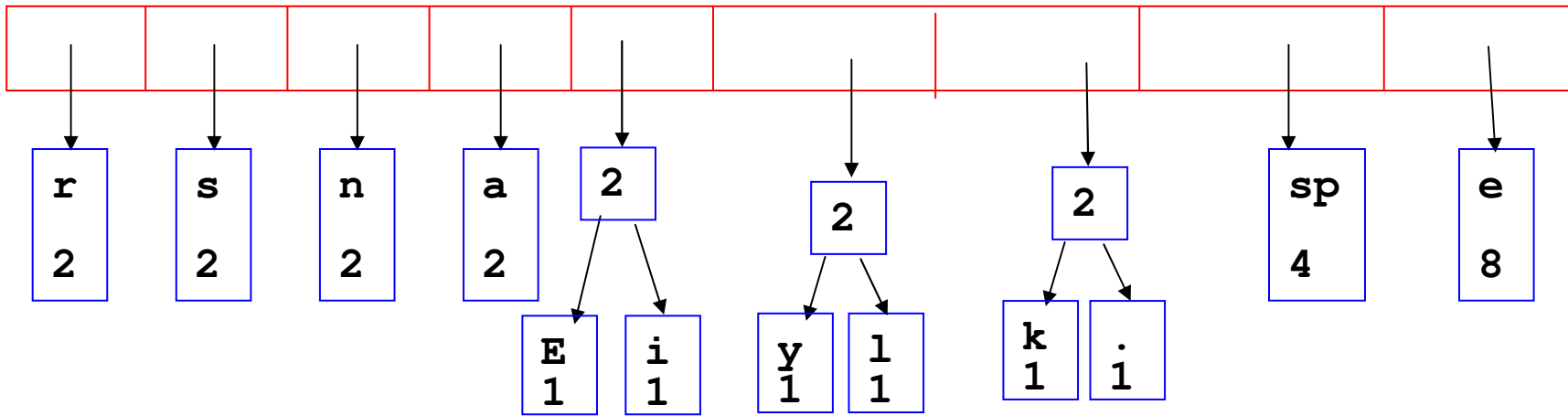
Building a Tree



Building a Tree

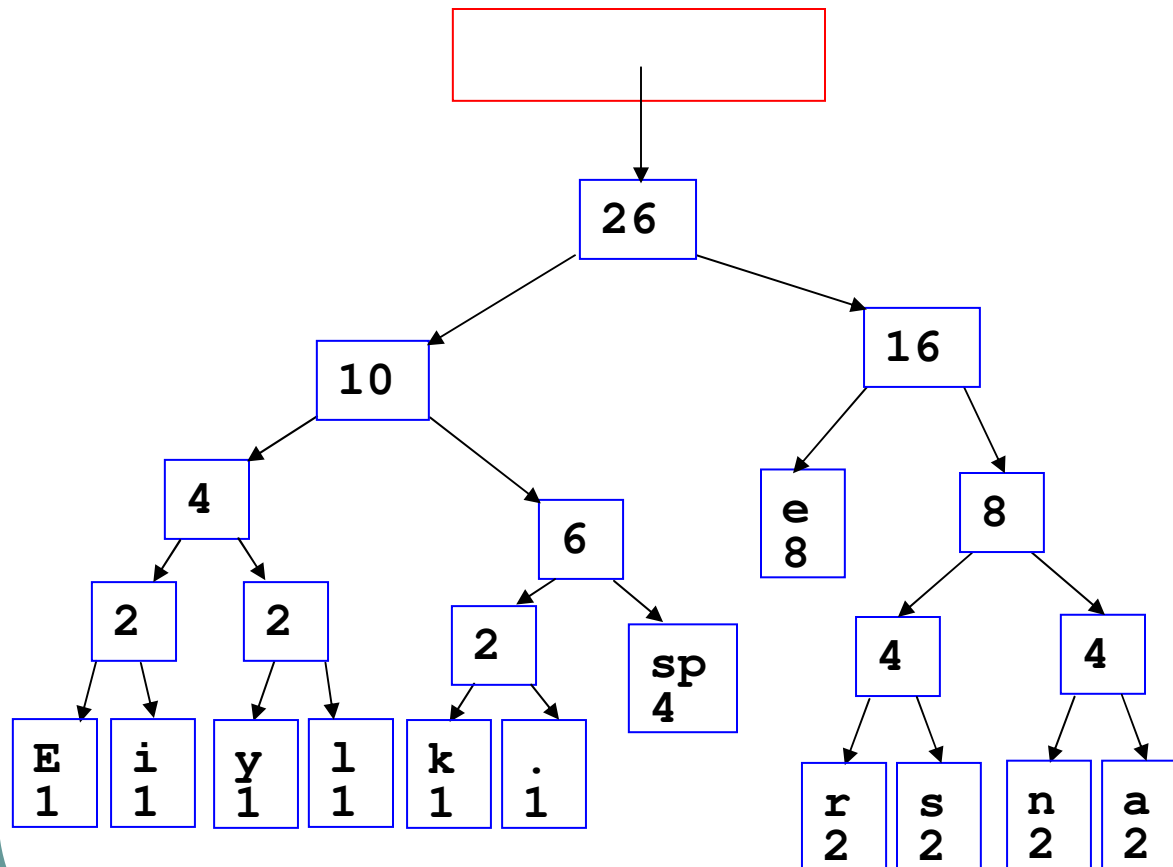


Building a Tree



- To continue ...

At the end



After enqueueing this node there is only one node left in priority queue.

How to implement ?

- Reuse JRB to represent the tree
 - Each new node is created as a JRB node
 - The edges are directional from the parents to the children.
 - Two edges are created and marked using label 0 or 1 when a parent node is created.
- Reuse Dlist or JRB to represent the priority queue
 - A queue node contains a key as the frequency of the related node in the tree
 - The queue node's value is a pointer referencing to the node in the tree

Quiz 1

- Reuse the graph API defined in previous class to write a function that builds a Huffman tree from a string as the following

```
typedef struct {  
    Graph graph;  
    JRB root;  
} HuffmanTree;  
HuffmanTree makeHuffman (char * buffer, int size);
```

Huffman code table

- In order to compress the data string, we have to build a code table from the Huffman tree. The following data structure is used to represent the code table

```
typedef struct {  
    int size;  
    char bits[2];  
} Coding;  
Coding huffmanTable[256];
```

- `huffmanTable['A']` give the coding of 'A'. If the coding's size = 0, the character 'A' is not present in the text. `bits` contains the huffman code (sequence of bits) of the given character.

Quiz 2

- Write a function to create the Huffman code table from a Huffman tree
 - `void createHuffmanTable(HuffmanTree htree, Coding*htable);`
- Write a function to compress a text buffer to a Huffman sequence.
 - `void compress(char * buffer, int size, char* huffman, int* nbit);`
- The buffer contains *size* characters. After compressing, the huffman buffer contains *nbit* bits for output.
- In order to write this function, you should create a function to add a new character into the huffman buffer as the following
 - `void addHuffmanChar(char * ch, Coding*htable, char* huffman, int* nbit);`