## Artificial Intelligence

First Order Logic

- Syntax
- Semantic

Lecturer 10 - First Order Logic

- Inference
- Resolution

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First Order Logic (FOL)

- First Order Logic is about
- Objects (e.g., people, houses)
- Relations(e.g., red, bigger than, father of)
- Facts
- The world is made of objects
- Objects are things with individual identities and properties to distinguish them
- Various relations hold among objects. Some of these relations are functional
- Every fact involving objects and their relations are either true or false


## FOL

- Syntax
- Semantic
- Inference
- Resolution


## FOL Syntax

## - Symbols

- Variables: $x, y, z, \ldots$
- Constants: a, b, c, ...
- Function symbols (with arities): f, g, h, ..
- Relation symbols (with arities): p, r, s
- Logical connectives: $\neg, \wedge, \vee, \Rightarrow, \Leftrightarrow$
- Quantifiers: $\quad \exists, \forall$


## FOL Syntax

- Variables, constants and function symbols are used to build terms
- X, Bill, FatherOf(X), ..
- Relations and terms are used to build predicates
- Tall(FatherOf(Bill)), Odd(X), Married(Tom,Marry), Loves(Y,MotherOf(Y)), ..
- Predicates and logical connective are used to build sentences
- Even(4)
- $\forall X$. Even $(X) \Rightarrow \operatorname{Odd}(X+1)$
- $\exists \mathrm{X} . \mathrm{X}>0$


## FOL

## FOL Semantic

- Syntax
- Variables
- Objects
- Constants
- Entities
- Function symbol

Function from objects to objects

- Relation symbol
- Relation between objects
- Quantifiers
- $\exists x . P$ true if P is true under some value of x
- $\forall x . P$ true if P is true under every value of x
- Logical connectives
- Similar to Propositional Logic


## Example

- Symbols
- Variables: x,y,z, ...
- Constants: $0,1,2, \ldots$
- Function symbols: +,*
- Relation symbols: >, =
- Semantic
- Universe: N (natural numbers)
- The meaning of symbols
- Constants: the meaning of 0 is the number zero, .
- Function symbols: the meaning of + is the natural number addition,
- Relation symbols: the meaning of $>$ is the relation greater than,..


## Robinson's Resolution for FOL

Given $K B=\{P 1(\ldots), \mathrm{P} 2(\ldots), \ldots, \mathrm{Pn}(\ldots)\}$. Prove $\mathrm{Q}(\ldots)$.
Add $\neg Q(\ldots)$ to $K B: K B=K B \wedge \neg Q(\ldots)$. Prove unsatisfied.
Theorem: A set of clauses $S$ is unsatisfiable if and only if upon the input $S$, Resolution procedure finds the empty clause (after a finite time).

1. Write each $\mathrm{Pi}(\ldots), \neg \mathrm{Q}(\ldots)$ in one line.
2. Transfer to CNF representation
$\forall x_{1} \forall x_{2} \ldots \forall x_{n}\left[p_{1}(\ldots) \vee \ldots \vee p_{n}(\ldots)\right] \wedge\left[q_{1}(\ldots) \vee \ldots \vee q_{m}(\ldots)\right] \quad(*)$
3. Break (*) into smaller clauses at the logic connective $\wedge$ :
$\forall x_{1} \forall x_{2} \ldots \forall x_{n}\left[p_{1}(\ldots) \vee \ldots \vee p_{n}(\ldots)\right]$
$\forall x_{1} \forall x_{2} \ldots \forall x_{n}\left[q_{1}(\ldots) \vee \ldots \vee q_{m}(\ldots)\right]$

## Transform sentences to FOL

|  | b. Anyone who owns a dog is a lover-of-animals |
| :---: | :---: |
|  | $\forall x .(\exists \mathrm{y} . \mathrm{D}(\mathrm{y}) \wedge(\mathrm{O}(\mathrm{x}, \mathrm{y})) \rightarrow \mathrm{L}(\mathrm{x})$ |
| a. John owns a dog | $\forall x \cdot(-\exists y \cdot(D(y) \wedge O(x, y)) \vee L(x)$ |
| $\exists x . D(x) \wedge O(J, x)$ | $\forall x . y y . \neg(D(y) \wedge O(x, y)) \vee L(x)$ |
| $D($ Fido $) \wedge O(J$, Fido $)$ | $\forall x . y y . \neg D(y) \vee \neg O(x, y) \vee L(x)$ |
|  | $\neg D(y) \vee \neg \mathrm{O}(\mathrm{x}, \mathrm{y}) \vee \mathrm{L}(\mathrm{x})$ |


| c. Lovers-of-animals do not kill <br> animals |
| :--- |
| $\forall x . L(x) \rightarrow(\forall y . A(y) \rightarrow-K(x, y))$ |
| $\forall x . \neg L(x) \vee(\forall y . A(y) \rightarrow \neg K(x, y))$ |
| $\forall x . \neg L(x) \vee(\forall y . \neg A(y) \vee \neg K(x, y))$ |
| $\neg L(x) v \neg A(y) \vee \neg K(x, y)$ |

## Robinson's Resolution for FOL

4. Resolution:
u) $\neg \mathrm{p}\left(\mathrm{x}_{1}, \mathrm{x}_{2}, \ldots, \mathrm{x}_{\mathrm{n}}\right) \vee \mathrm{q}(\ldots)$
v) $p\left(y_{1}, y_{2}, \ldots, y_{n}\right) \vee r(\ldots)$
with substitution $\theta=\left\{\frac{z_{1}}{x_{1}}, \frac{z_{1}}{y_{1}}, \ldots, \frac{z_{n}}{x_{n}}, \frac{z_{n}}{y_{n}}\right\}$
5. Contrast appears when KB contains 2 lines:
i) $\neg p\left(x_{1}, x_{2}, \ldots, x_{n}\right) \quad$ ii) $\left.\left.p\left(y_{1}, y_{2}, \ldots, y_{n}\right)\right\} \Rightarrow w\right) q(\ldots) \vee r(\ldots)$ with substitution

$$
\theta=\left\{\frac{z_{1}}{x_{1}}, \frac{z_{1}}{y_{1}}, \ldots, \frac{z_{n}}{x_{n}}, \frac{z_{n}}{y_{n}}\right\}
$$

## Step 4 - Example

- Which substitution $\theta$ is used for resolving:
$P(a, x, b)$, and
$\neg P(y, z, z)$

$$
\theta=\left\{\frac{a}{y}, \frac{b}{z}, \frac{b}{x}\right\}
$$

- $P(a, b, b)$
- ${ }_{\neg} P(a, b, b)$


## Step 4 - Example

- Which substitution $\theta$ is used for resolving:


## Step 4 - Example

$$
\frac{\neg \operatorname{Pet}(J o e) \vee \operatorname{Cat}(J o e) \vee \operatorname{Bird}(J o e) \quad \operatorname{Parrot}(x) \vee \neg \operatorname{Bird}(x)}{\neg \operatorname{Pet}\left({ }^{\frac{1}{2}}\right)} \text { (1) }
$$

(1) $m g u(\operatorname{Bird}(x), \operatorname{Bird}(J o e))=\{x / J o e\}$

$$
\begin{aligned}
& \quad \frac{\neg \operatorname{On}(x, y) \vee \operatorname{Above}(x, y) \quad O n(B, A) \vee \operatorname{On}(A, B)}{\operatorname{Above}(A, B) \vee O n(B, A)} \text { (2) } \\
& \text { (2) } m u g u(O n(x, y), O n(A, B))=\{x / A, y / B\}
\end{aligned}
$$

$$
\frac{\neg \operatorname{Bird}(x) \vee \text { Feathers }(x) \quad \neg \text { Feathers }(y) \vee \text { Flies }(y)}{\neg \operatorname{Bird}(x) \vee \text { Flies }(x)}
$$

(3) $m g u($ Feathers $(x)$, Feathers $(y))=\{y / x\}$

## Step 4 - Example

- Given $p(a, b), p(c, d), q(d, c, c)$ are true
- Rule
$p(x, y) \wedge q(y, x, x) \Rightarrow r(x, y)$
- Using substitutions with the above rule, generate new predicate
- Idea:
- Try with $p(x, y) \equiv p(a, b)$ or $p(x, y) \equiv p(c, d)$


## Resolution - Example

$$
\begin{array}{cc}
\forall x & P(x) \rightarrow Q(x) \\
\forall x & \neg P(x) \rightarrow R(x)
\end{array} \quad \text { Resolve } 1 \text { and } 30 . \neg P(x) \vee S(x)
$$

## Exercice 1

- John owns a dog
- Anyone who owns a dog is a lover of animals
- Lovers of animals do not kill animals
- Proves that John does not kill animals?

Transform the problem to a set of clauses and apply Robinson's resolution

| $\exists x . \mathrm{D}(\mathrm{x}) \wedge \mathrm{O}($ John, x$)$ | $\mathrm{D}($ Fido $) \wedge \mathrm{O}($ John, Fido $)$ |
| :--- | :--- |
| $\forall \mathrm{x} .(\exists \mathrm{y} . \mathrm{D}(\mathrm{y}) \wedge \mathrm{O}(\mathrm{x}, \mathrm{y})) \rightarrow \mathrm{L}(\mathrm{x})$ | $\neg \mathrm{D}(\mathrm{y}) \vee \neg \mathrm{O}(\mathrm{x}, \mathrm{y}) \vee \mathrm{L}(\mathrm{x})$ |
| $\forall \mathrm{x} . \mathrm{L}(\mathrm{x}) \rightarrow(\forall \mathrm{y} . \mathrm{A}(\mathrm{y}) \rightarrow \neg \mathrm{K}(\mathrm{x}, \mathrm{y}))$ | $\neg \mathrm{L}(\mathrm{x}) \vee \neg \mathrm{A}(\mathrm{y}) \vee \neg \mathrm{K}(\mathrm{x}, \mathrm{y})$ |
| $\forall \mathrm{x} . \mathrm{D}(\mathrm{x}) \Rightarrow \mathrm{A}(\mathrm{x})$ | $\neg \mathrm{D}(\mathrm{x}) \vee \mathrm{A}(\mathrm{x})$ |
| $\forall \mathrm{x} . \mathrm{A}(\mathrm{x}) \Rightarrow \neg \mathrm{K}($ John, x$)$ | $\mathrm{A}($ Fido $) \wedge \mathrm{K}($ John, Fido $)$ |

## Exercice 2

- Jack owns a dog
- Every dog owner is an animal lover
- No animal lover kills an animal
- Either Jack or Curiosity killed the cat, who is named Tuna
- Prove that Curiosity kill the cat.


## Jack owns a dog own(Jack, dog) <br> Every dog owner is an animal lover <br> No animal lover kills an animal <br> Either Jack or Curiosity killed the cat, who is named Tuna <br> Did Curiosity kill the cat? Kills(Curiosity,Tuna)

$\exists x . \operatorname{Dog}(x) \wedge O w n s($ Jack, $x)$
$\forall x \forall y .(\operatorname{Dog}(y) \wedge O w n s(x, y)) \Rightarrow$ AnimalLover $(x)$
$\forall x .(\exists y . \operatorname{Dog}(y) \wedge O w n s(x, y)) \Rightarrow$ AnimalLover $(x)$
$\forall x \forall y .($ AnimalLover $(x) \wedge \operatorname{Animal}(y) \Rightarrow \neg \operatorname{Kills}(x, y))$
Kills(Jack,Tuna) $\vee \operatorname{Kill}($ Curiosity,Tuna)
Cat(Tuna)
$\forall x \cdot \operatorname{Cat}(x) \Rightarrow \operatorname{Animal}(x)$

Transform the problem to set of clauses
$\operatorname{Dog}(D)$
Owns(Jack, D)
$\neg \operatorname{Dog}(y) \vee \neg$ Owns $(x, y) \vee$ AnimalLover $(x)$
$\neg$ AnimalLover $(x) \wedge \neg$ Animal $(y) \vee \neg \operatorname{Kills}(x, y)$
Kills(Jack,Tuna) $\vee \operatorname{Kill}(C u r i o s i t y, T u n a)$
Cat(Tuna)
$\neg \operatorname{Cat}(x) \vee \operatorname{Animal}(x)$
$\neg$ Kills(Curiosity,Tuna)

## Exercice 3

- The law says that it is a crime for an American to sell weapons to hostile nations
- The country Nono, an enemy of America, has some missiles, and all of its missiles were sold to it by Colonel West, who is American
- Is West a criminal?

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The country Nono, an enemy of America, has some missiles, and all of its missiles were
sold to l= by Golonel West, who is American
sold to \jby Golonel West
    "... it is a crime for an American to sell weapons to hostile nations";
        V }x,y,z American(x) A Weapon(y)A Nation(z)A Hostile(z
        A Selis(x, z, y) = Criminal(x)
    "Nono ... has some missiles"
        \existsx Owns(Nono,x)^Missile(x)
    "All of its missiles were sold to it by Colonel West"
        \forallx Owns(Nono,x)A Missile(x) => Sells(West.Nono,x)
    We will also need to know that missiles are weapons
        \forall Missile(x) => Weapon(x)
    and that an enemy of America counts as "hostile"
        \forallx Enemy(x.America) }=>\mathrm{ Hostile(x)
    "West, who is American ...":
        American(West)
    "The country Nono ...":
        Nation(Nono)
            "Nono, an enemy of America ...":
        Enemy(Nono, America)
        Nation(America)

Transform the problem to set of clauses and Resolution
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