

The background features a white surface with scattered, colorful abstract shapes. These include thin, curved lines in shades of light blue, light green, and light purple. Interspersed among these are numerous small, yellow, triangular shapes that resemble confetti or paper scraps. The overall aesthetic is clean and modern.

C Programming Basic – week 8

The slide features a decorative background on the left side with a light green balloon at the top, a light blue balloon in the middle, and a light purple balloon at the bottom. Yellow streamers and triangular flags are scattered around the balloons. The word "Content" is centered in a large, bold, dark teal font.

Content

1. How to use debugger tool (gdb)
2. Binary tree
3. Binary search tree
4. Recursive processing on Tree

1. gdb for debugging (1)

- **gdb**: the Gnu DeBugger
- <http://www.cs.caltech.edu/courses/cs11/material/c/mike/misc/gdb.html>
- Use when program core dumps
- or when want to walk through execution of program line-by-line

`gdb` for debugging (2)

- Before using `gdb`:
 - Must compile C code with additional flag:
`-g`
 - This puts all the source code into the binary executable
- Then can execute as: `gdb myprogram`
- Brings up an interpreted environment

`gdb` for debugging (3)

`gdb> run`

- Program runs...
- If all is well, program exits successfully, returning you to prompt
- If there is (*e.g.*) a core dump, `gdb` will tell you and abort the program

gdb – basic commands (1)

- Stack backtrace ("**where**")
 - Your program core dumps
 - Where was the last line in the program that was executed before the core dump?
 - That's what the **where** command tells you

gdb - basic commands (2)

gdb> where

last call

last call in your code

```
#0 0x4006cb26 in free () from /lib/libc.so.6
```

```
#1 0x4006ca0d in free () from /lib/libc.so.6
```

```
#2 0x8048951 in board_updater (array=0x8049bd0,  
ncells=2) at 1dCA2.c:148
```

```
#3 0x80486be in main (argc=3, argv=0xbffff7b4) at  
1dCA2.c:44
```

```
#4 0x40035a52 in __libc_start_main () from  
/lib/libc.so.6
```

stack backtrace

gdb – basic commands (3)

- Look for topmost location in stack backtrace that corresponds to your code
- Watch out for
 - freeing memory you didn't allocate
 - accessing arrays beyond their maximum elements
 - dereferencing pointers that don't point to part of a `malloc()` ed block

gdb – basic commands (4)

- **break, continue, next, step** commands
- **break** causes execution to stop on a given line

```
gdb> break foo.c: 100    (setting a  
breakpoint)
```

- **continue** resumes execution from that point
- **next** executes the next line, then stops
- **step** executes the next statement
 - goes into functions if necessary (**next** doesn't)

gdb - basic commands (5)

- **print** and **display** commands
- **print** prints the value of any program expression

```
gdb> print i
```

```
$1 = 100
```

- **display** prints a particular value every time execution stops

```
gdb> display i
```

`gdb` – printing arrays (1)

- `print` will print arrays as well

```
int arr[] = { 1, 2, 3 };
```

```
gdb> print arr
```

```
$1 = {1, 2, 3}
```

- N.B. the `$1` is just a name for the result

```
print $1
```

```
$2 = {1, 2, 3}
```

`gdb` - printing arrays (2)

- `print` has problems with dynamically-allocated arrays

```
int *arr;
```

```
arr = (int *)malloc(3 * sizeof(int));
```

```
arr[0] = 1; arr[1] = 2; arr[2] = 3;
```

```
gdb> print arr
```

```
$1 = (int *) 0x8094610
```

- Not very useful...

gdb - printing arrays (3)

- Can print this array by using @ (gdb special syntax)

```
int *arr;
```

```
arr = (int *)malloc(3 * sizeof(int));
```

```
arr[0] = 1; arr[1] = 2; arr[2] = 3;
```

```
gdb> print *arr@3
```

```
$2 = {1, 2, 3}
```

`gdb` – abbreviations

- Common `gdb` commands have abbreviations
 - `p` (same as `print`)
 - `c` (same as `continue`)
 - `n` (same as `next`)
 - `s` (same as `step`)
- More convenient to use when interactively debugging

other instruction

- clear : delete break point of current file.
- delete [break position]: delete breakpoint at a specific file and position
- Conditional break

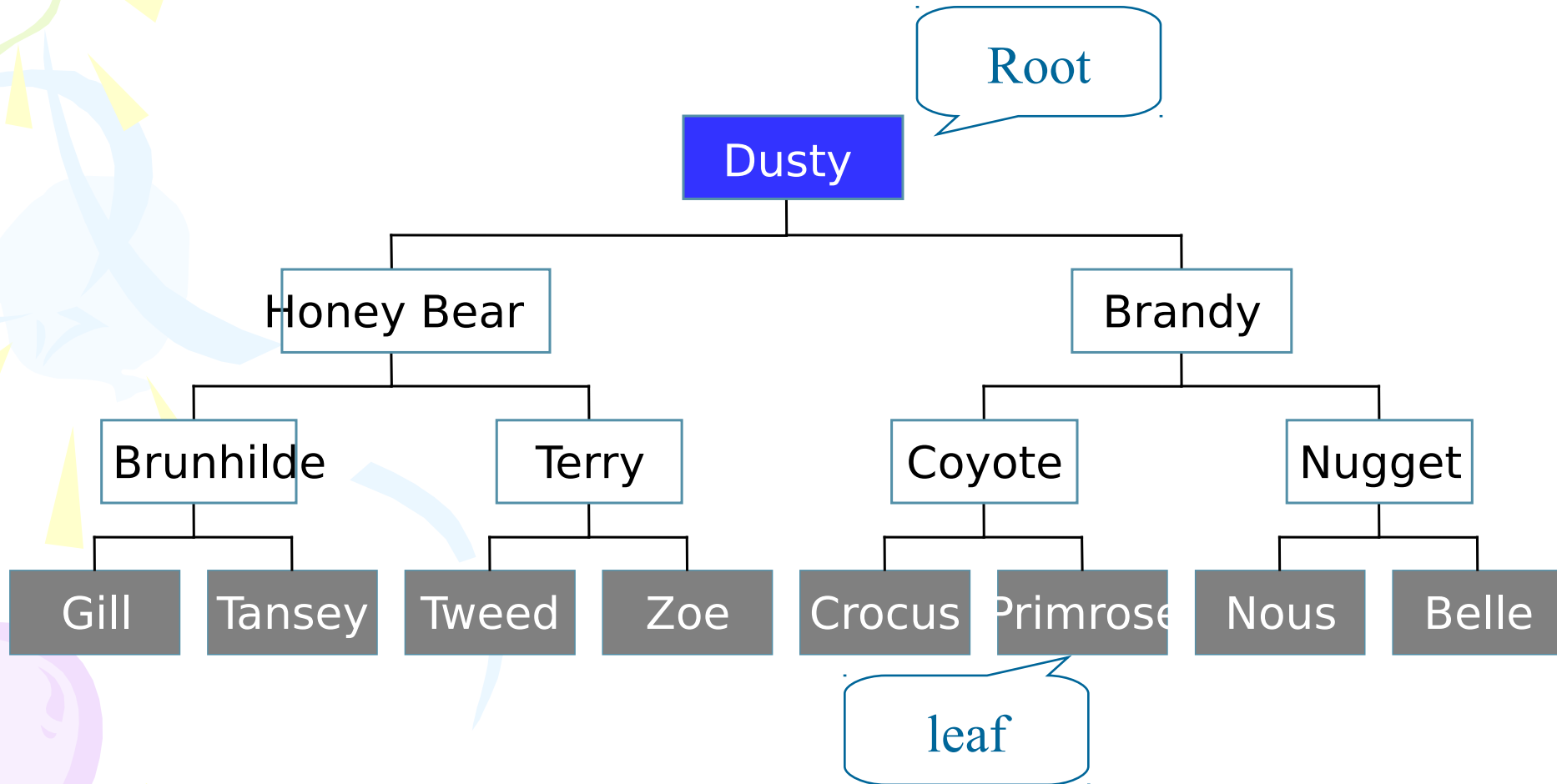
```
gdb> break foo.c: 100 if i==-1
```

- quit
- run: restart from beginning.

2. Binary tree

- Linked lists are **linear structures** and it is difficult to use them to organize an **hierarchical** representation of objects
- Although stacks and queues reflect some hierarchy, they are limited to only **one dimension**
- To overcome this limitation, we create a new data type called a **tree** that consists of **nodes** and **arcs**. Unlike natural trees, these trees are **depicted upside down** with the **root** at the top and the **leaves** at the bottom.

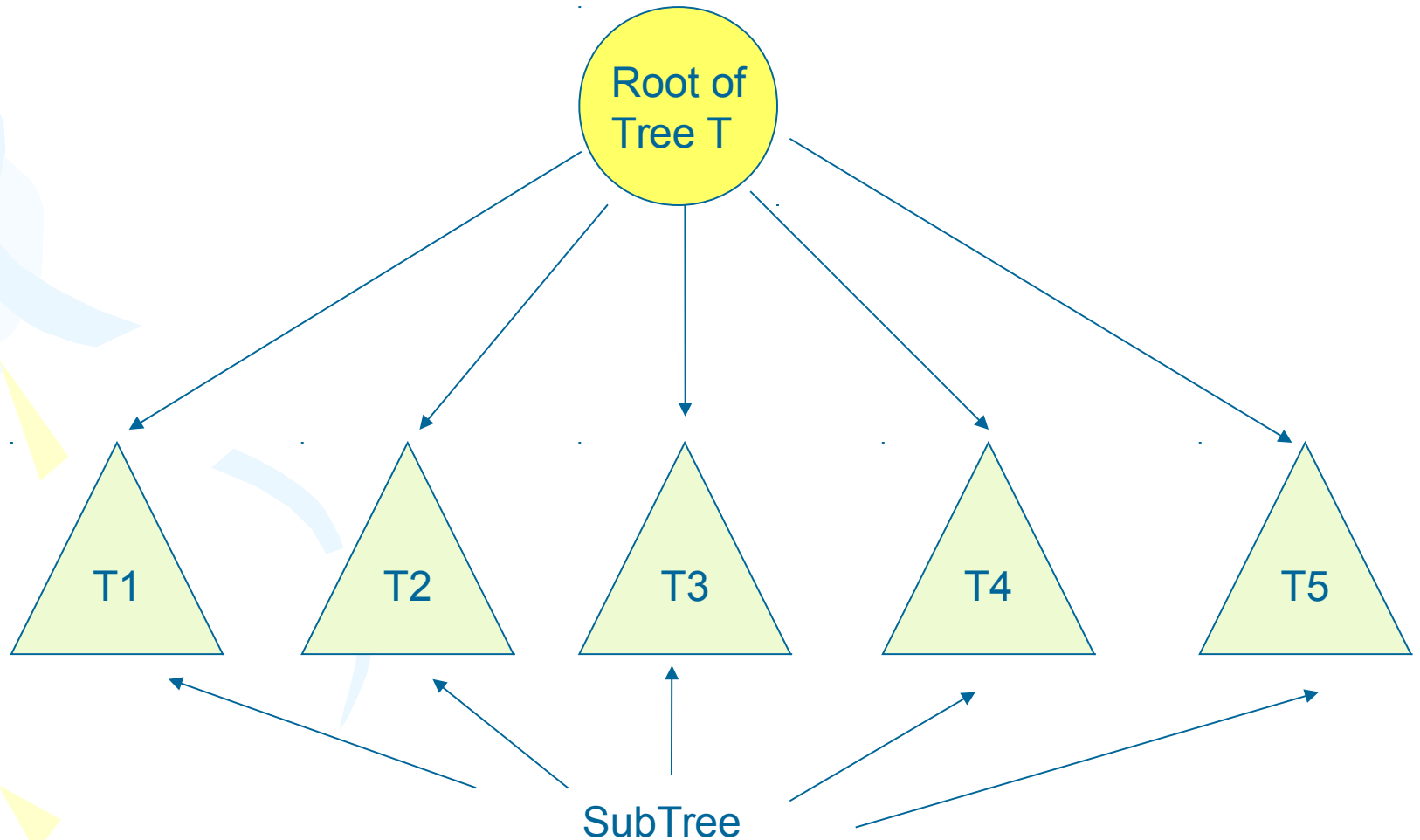
Family Tree



Definition of tree

- A tree is a finite set of one or more nodes such that:
- There is a specially designated node called the root.
- The remaining nodes are partitioned into $n \geq 0$ disjoint sets T_1, \dots, T_n , where each of these sets is a tree.
- We call T_1, \dots, T_n the subtrees of the root.

Recursive definition



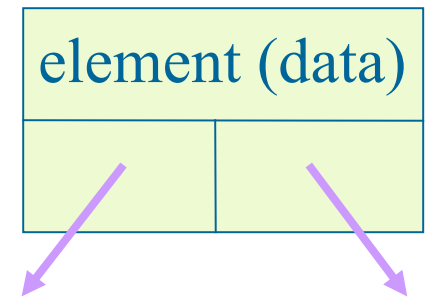
Binary Tree

- A binary tree is a tree in which no node can have more than two children.
- Each node has 0, 1, or 2 children

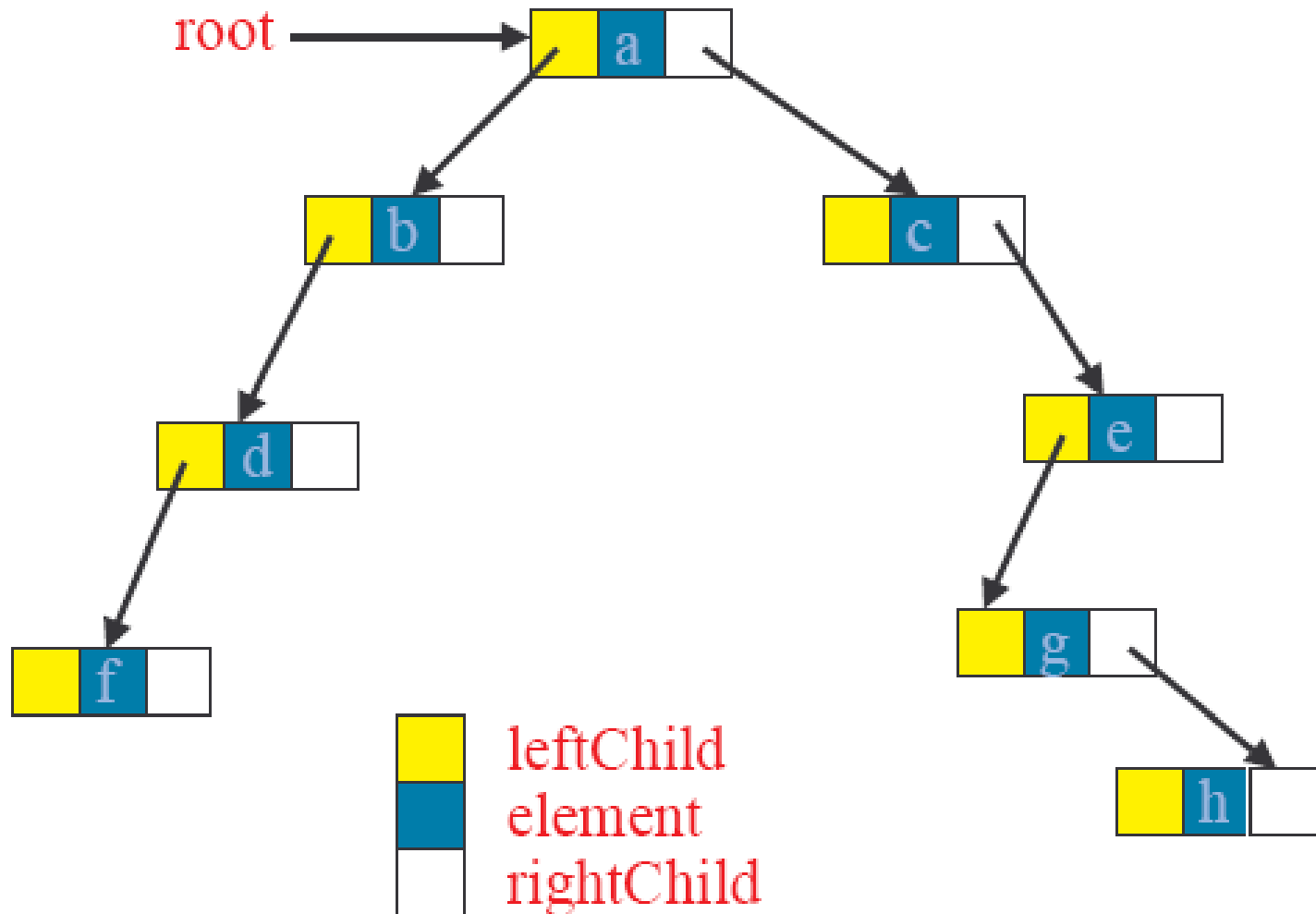
Linked Representation

- Each tree node is represented as an object whose data type is
- The space required by an n node binary tree is $n *$ (space required by one node)

```
typedef ... elmType;
//whatever type of element
typedef struct nodeType {
    elmType element;
    struct nodeType *left, *right;
};
typedef struct nodeType *treeType;
```



A linked binary tree



Binary Tree ADT

- `makeNullTree(treeType *t)`
- `creatNewNode()`
- `isEmpty()`

Tree initialization and verification

```
typedef ... elmType;
typedef struct nodeType {
    elmType element;
    struct nodeType *left, *right;
} node_Type;

typedef struct nodeType *treeType;

void makeNullTree(treeType *T) {
    (*T)=NULL;
}

int emptyTree(treeType T) {
    return T==NULL;
}
```


Access left and right child

```
treeType leftChild(treeType n)
{
    if (n!=NULL) return n->left;
    else return NULL;
}
treeType rightChild(treeType n)
{
    if (n!=NULL) return n->right;
    else return NULL;
}
```

Create a new node

```
nodeType *createNode (elmType NewData)
{
    nodeType *N;
    N=(nodeType*)malloc (sizeof (nodeType) );
    if (N != NULL)
    {
        N->left = NULL;
        N->right = NULL;
        N->element = NewData;
    }
    return N;
}
```

Check if a node is a leaf

```
int isLeaf(treeType n) {  
    if (n!=NULL)  
        return (leftChild(n)==NULL) &&  
            (rightChild(n)==NULL);  
    else return -1;  
}
```

Recursive processing: Number of nodes

- As tree is a recursive data structure, recursive algorithms are useful when they are applied on tree.

```
int nb_nodes(treetype T) {  
    if (EmptyTree(T)) return 0;  
    else return 1+nb_nodes(LeftChild(T)) +  
        nb_nodes(RightChild(T));  
}
```

Create a tree from two subtrees

```
treetype createfrom2 (elmttype v,  
    treetype l, treetype r) {  
    treetype N;  
    N=(node_type*)malloc(sizeof(node_type));  
    N->element=v;  
    N->left=l;  
    N->right=r;  
    return N;  
}
```

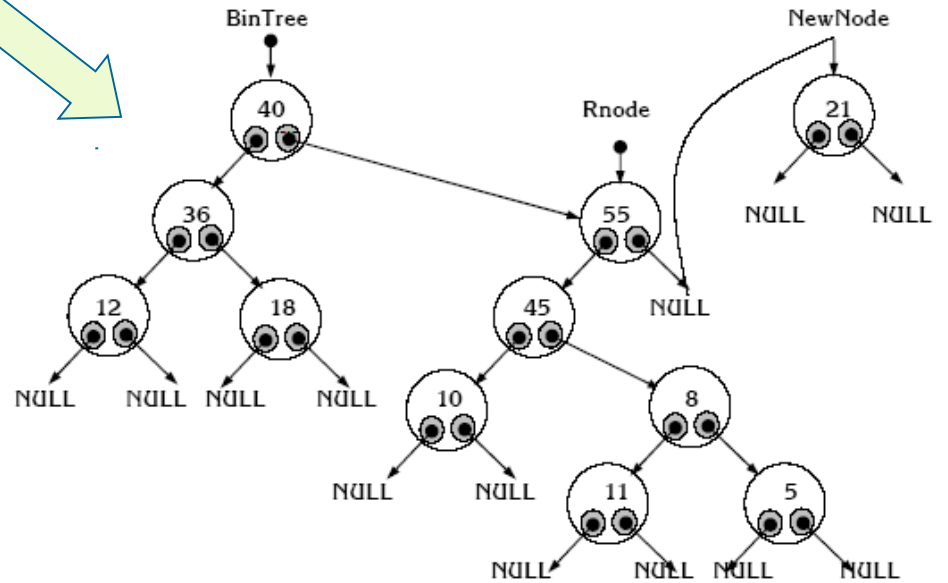
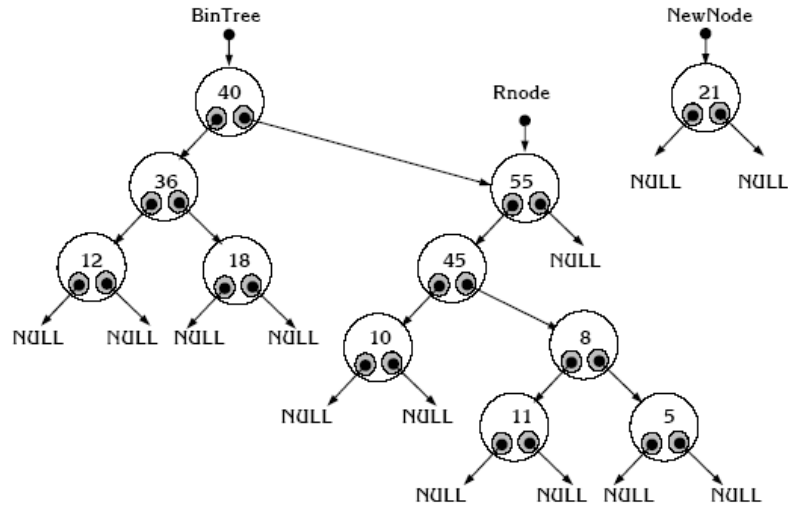
Adding a new node to the left most position

```
treetype Add_Left(treetype *Tree, elmtyp NewData)
{ node_type *NewNode = Create_Node(NewData);
  if (NewNode == NULL) return (NewNode);
  if (*Tree == NULL)
    *Tree = NewNode;
  else{
    node_type *Lnode = *Tree;
    while (Lnode->left != NULL)
      Lnode = Lnode->left;
    Lnode->left = NewNode;
  }
  return (NewNode);
}
```

Adding a new node to the right most position

```
treetype Add_Left(treetype *Tree, elmtyp NewData)
{ node_type *NewNode = Create_Node(NewData);
  if (NewNode == NULL) return (NewNode);
  if (*Tree == NULL)
    *Tree = NewNode;
  else{
    node_type *Rnode = *Tree;
    while (Rnode->right != NULL)
      Rnode = Rnode->right;
    Rnode->right = NewNode;
  }
  return (NewNode);
}
```

Illustration

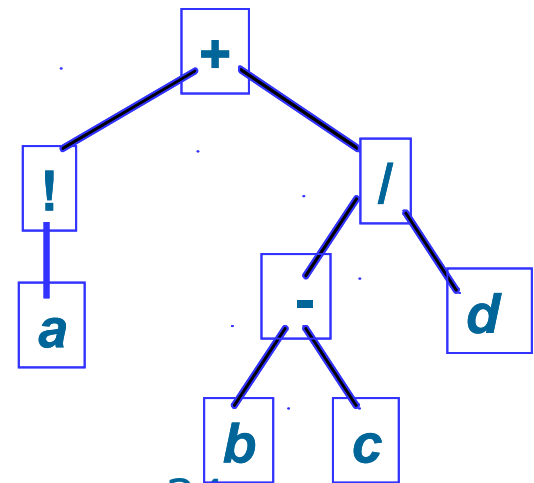


Homework 1

- Develop the following helper functions for a tree:
 - return the height of a binary tree.
 - return the number of leafs
 - return the number of internal nodes
 - count the number of right children.

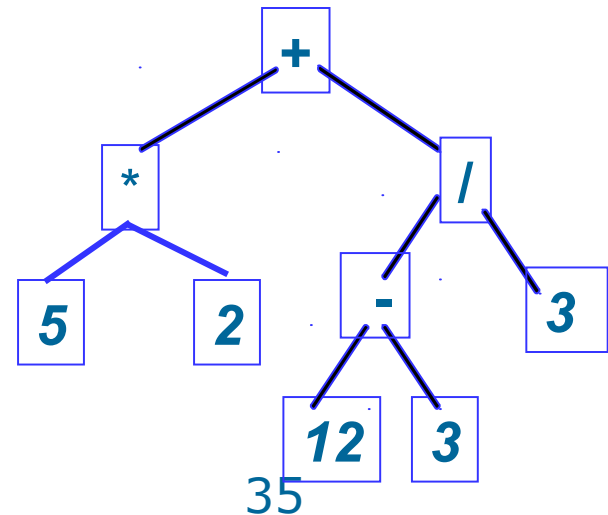
Exercise 8.1

- A binary tree can represent an arithmetic expression: The leaves are operands and the other nodes are operators.
- The left and right subtrees of an operator node represent **subexpressions** that must be evaluated **before** applying the operator at the root of the subtree.
- For example
 $a + (b - c)/d$
- Write a program to create a tree representing this expression



Homework 2

- Write an menu program that take a valid arithmetic expression as input and:
 - Store and represent it in a tree
 - Evaluate the expression.

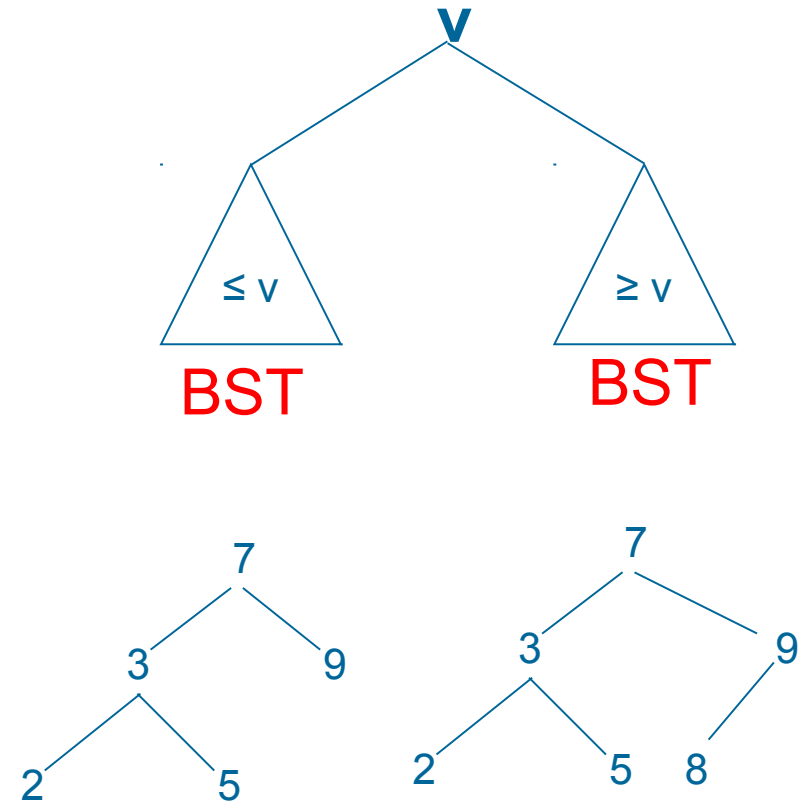


Homework 3

- Create file USopen.txt with the following content
 - Line 1 contains 16 players
- Build binary tree representation match results until the final match. At first, 16 players are leaf nodes.
- Choose randomly a winner in a match.
- Print result to file treegame.txt

3. Binary Search Tree

- Every element has a unique key.
- The keys in a nonempty left subtree (right subtree) are smaller (larger) than the key in the root of subtree.
- The left and right subtrees are also binary search trees.



Binary Search Tree Implementation

```
#include <stdio.h>
#include <stdlib.h>
typedef . . . KeyType; // specify a type
for the data
typedef struct Node{
    KeyType key;
    struct Node* left, right;
} NodeType;
typedef Node* TreeType;
```

Search on BST

```
TreeType Search(KeyType x, TreeType Root) {
    if (Root == NULL) return NULL; // not found
    else if (Root->key == x) /* found x */
        return Root;
    else if (Root->key < x)
        //continue searching in the right sub tree
        return Search(x, Root->right);
    else {
        // continue searching in the left sub tree
        return Search(x, Root->left);
    }
}
```

Insert a node to BST

- In a binary, there are not two nodes with the same key.

```
void InsertNode(KeyType x, TreeType *Root ) {
    if (*Root == NULL) {
        /* Create a new node for key x */
        *Root = (NodeType*) malloc (sizeof (NodeType));
        (*Root) -> key = x;
        (*Root) -> left = NULL;
        (*Root) -> right = NULL;
    }
    else if (x < (*Root) -> key) InsertNode(x,
        &(*Root) -> left);
    else if (x > (*Root) -> key) InsertNode(x, &(*Root) -
        > right);
}
```

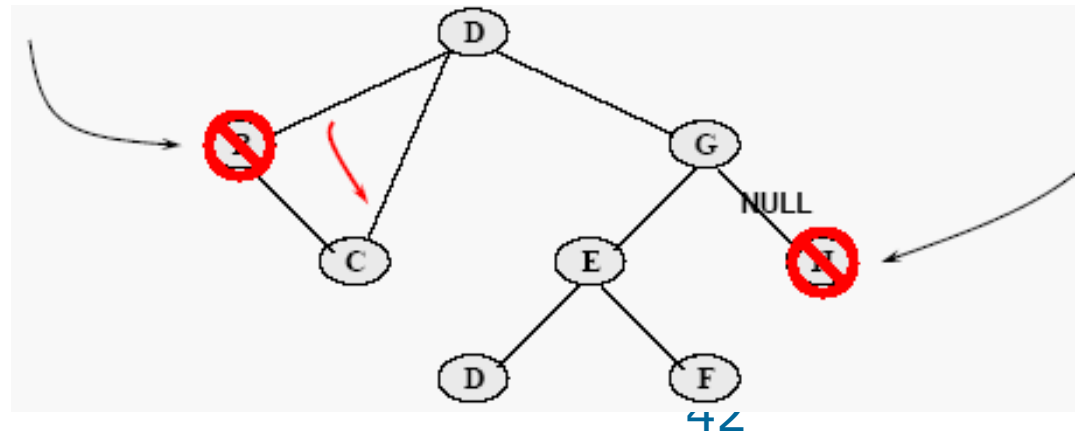

Insert a node to BST

- Version with the return type

```
TreeType InsertNode(KeyType x, TreeType Root ) {
if (Root == NULL) {
    /* Create a new node for key x */
    Root=(NodeType*)malloc(sizeof(NodeType));
    Root->key = x;
    Root->left = NULL;
    Root->right = NULL;
    Return Root;
}
else if (x < Root->key) return InsertNode(x, Root->left);
else if (x > Root->key) return InsertNode(x, Root->right);
}
```

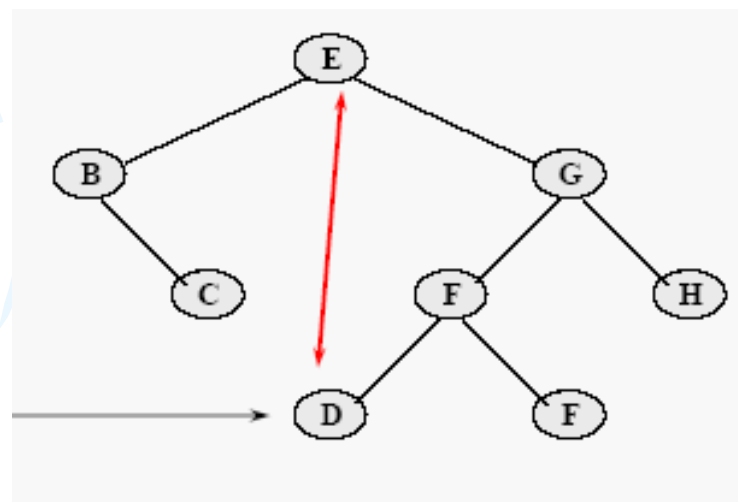
Delete a node from a BST

- Removing a leaf node is trivial, just set the relevant child pointer in the parent node to NULL.
- Removing an internal node which has only one subtree is also trivial, just set the relevant child pointer in the parent node to target the root of the subtree.



Delete a node from a BST

- Removing an internal node which has two subtrees is more complex
 - Find the left-most node of the right subtree, and then swap data values between it and the targeted node.
 - Delete the swapped value from the right subtree.



Find the left-most node of right sub tree

- This function find the leftmost node then delete it.

```
KeyType DeleteMin (TreeType *Root ) {
    KeyType k;
    if ((*Root)->left == NULL) {
        k=(*Root)->key;
        (*Root) = (*Root)->right;
        return k;
    }
    else return DeleteMin(&(*Root)->left);
}
```

Delete a node from a BST

```
void DeleteNode(key X, TreeType *Root) {  
    if (*Root != NULL)  
        if (x < (*Root)->Key) DeleteNode(x, &(*Root)-  
            >left)  
        else if (x > (*Root)->Key)  
            DeleteNode(x, &(*Root)->right)  
        else if  
            ((*Root)->left == NULL) && ((*Root)->right == NULL)  
            *Root = NULL;  
        else if ((*Root)->left == NULL)  
            *Root = (*Root)->right  
        else if ((*Root)->right == NULL)  
            *Root = (*Root)->left  
        else (*Root)->Key = DeleteMin(&(*Root)->right);  
}
```

Pretty print a BST

```
void prettyprint(TreeType tree, char *prefix) {
    char *prefixend=prefix+strlen(prefix);
    if (tree!=NULL) {
        printf("%04d", tree->key);
        if (tree->left!=NULL) if (tree->right==NULL) {
            printf("\304"); strcat(prefix, "    ");
        }
        else {
            printf("\302"); strcat(prefix, "\263    ");
        }
        prettyprint(tree->left, prefix);
        *prefixend='\0';
        if (tree->right!=NULL) if (tree->left!=NULL) {
            printf("\n%s", prefix); printf("\300");
        } else printf("\304");
        strcat(prefix, "    ");
        prettyprint(tree->right, prefix);
    }
}
```

Exercise 8.2

- Write a function to delete all node of a tree. This function must be called before terminating program.

Exercise 8.3

- Create a binary search tree with 10 nodes. Each node contains a random integer.
- Ask user to input a number and search for it.
- Print the content of the trees.

Exercise 8.4

- Write functions FindMin and FindMax for BST library
- Argument: root pointer
- Return: pointer to the min/max node.

Exercise 8.5

- We assume that you make a mobile phone's address book.
- Declare a structure which can store at least "name", "telephone number", "e-mail address."
- Read 10 addresses from an input file to a BST in increasing order of email address
- (1) Find a specified e-mail address print to a file if found.
- (2) Output all the data stored in the binary tree in ascending order for the e-mail address.

Homework 4

- Write an program for looking up English-Vietnamese information technology dictionary. Vietnamese without diacritic.
- Instruction: An entry contains English word – Vietnamese meaning.
- Use BST to store.
- Basic functions: search word, add word, delete word and write to file.
- Advance: Each word has a list of synonym

Homework 5

- Check search speed of a BST.
- Generate 1 million random integers and insert to BST.
- Print tree height.
- The program allows to:
 - Create a new tree (remember to free memory for the last tree).
 - Search and print out the number of comparisons

Homework 6

- Continue the program for NokiaDB.
- Search a mobile model using BST.
- Functions: Import, Insert, Delete, Update, Search, Print the list.

Instruction

- Separate data structure library (e.g. BST) and the main program